## **Tuning The Trimming Factor**

For arbitrary lower and upper bounds lb and ub, there are certain values of the trimming factor that can negatively impact correctness and query efficiency.<sup>1</sup> In this document, we describe two mappings that, given any domain, will output an appropriate trimming factor that maintains correctness, while achieving a good trade-off between query latency, insertion throughput, and storage overhead. In the following, we assume that there exist a default trimming factor  $tf \in \mathbb{N}_{\geq 0}$ as well as a default sparsity factor  $sp \in \{1, 2, 4, 8\}$ .<sup>2</sup>

The integer case. The mapping,  $map_I$ , applies to the following numerical data types: sint32, int64 and sint64 and is defined as follows.

$$\mathsf{map}_{I}(\mathsf{lb},\mathsf{ub}) = \begin{cases} \bot & \text{if } \mathsf{ub} < \mathsf{lb} \\ \min\left(\mathsf{tf}, \left\lfloor \mathsf{sp}^{-1} \cdot \left( \left\lceil \log(\mathsf{ub} - \mathsf{lb} + 1) \right\rceil \right) \right\rfloor \right) & \text{otherwise} \end{cases}$$

where lb and ub belong to the underlying domain of the numerical data type.

The floating-point format case. The mapping ,  $map_F$ , applies to the following numerical data types: bin128 and dec128 and is defined as follows.

$$\mathsf{map}_F(\mathsf{lb},\mathsf{ub},\mathsf{prc}) = \left\{ \begin{array}{ll} \bot & \text{if } \mathsf{ub} < \mathsf{lb} \\ \min\left(\mathsf{tf},\left\lfloor \mathsf{sp}^{-1} \cdot \left(\left\lceil \log((\mathsf{ub} - \mathsf{lb} + 1) \cdot 10^{\mathsf{prc}})\right\rceil\right)\right\rfloor\right) & \text{otherwise.} \end{array} \right.$$

where  $\mathsf{lb}$  and  $\mathsf{ub}$  belong to the underlying domain of the numerical data type and where  $\mathsf{prc} \in \mathbb{N}_{\geq 0}$ .

<sup>&</sup>lt;sup>1</sup>For example, if both the domain and the chosen trimming factor are large, then at query time, there are cases where the size of the cover can be extremely large. A very large cover has two implications: an efficiency implication as the number of binary hops is proportional to the size of the cover; and a correctness implication, as the number of generated tags in this case can be too high exceeding the BSON limit.

<sup>&</sup>lt;sup>2</sup>Our experiments show that setting the trimming factor to 6 and the sparsity factor to 2 achieves the best trade-off between query latency and insertion throughput.